

Panorama des modèles et outils de programmation parallèle

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Outline

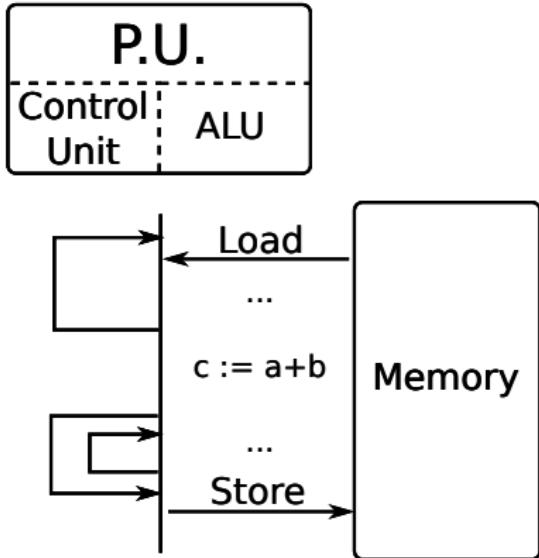
Introduction

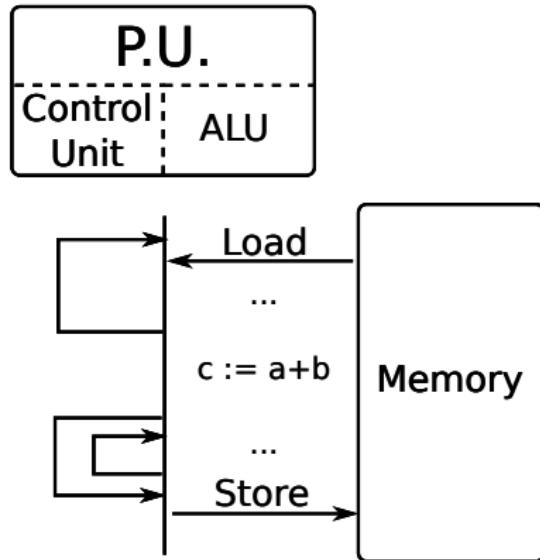
Accelerators & Many-Core

Runtime Systems

Conclusion

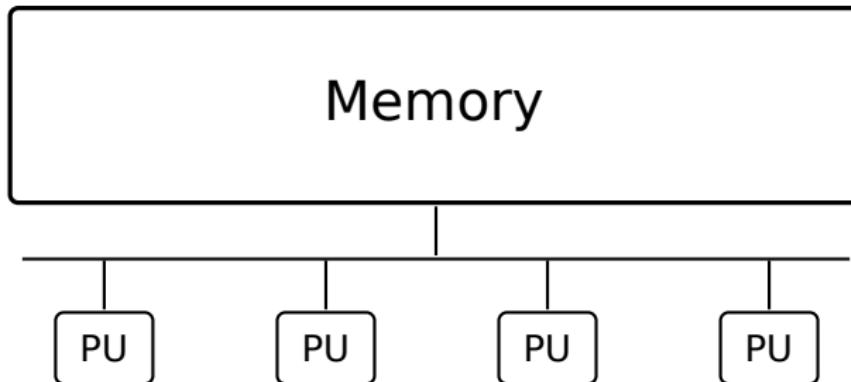
Introduction

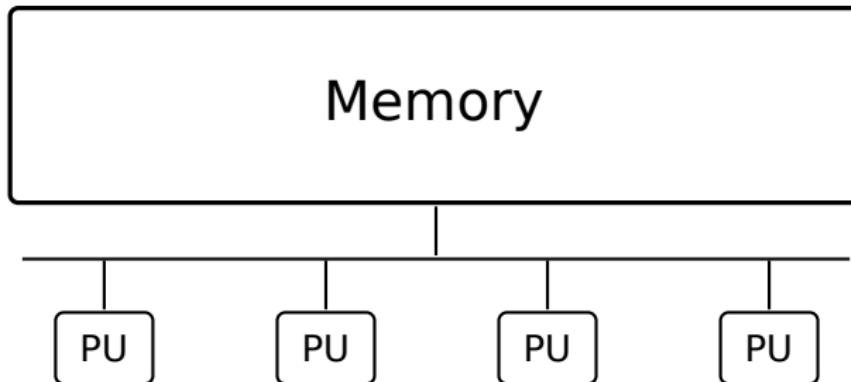




- ▶ asm (x86...), C, Fortran, C++...

Moore's Law + Frequency Wall = Parallelism
(mandatory slide)





- ▶ PThread

POSIX Threads (PThread)

```
#include <pthread.h>

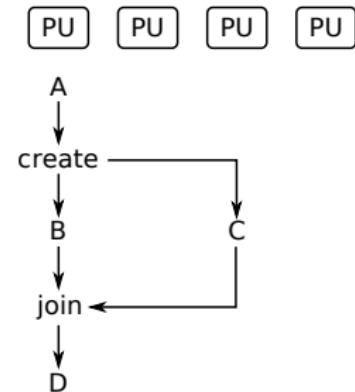
void * mythread(void * arg) {
    // C
}

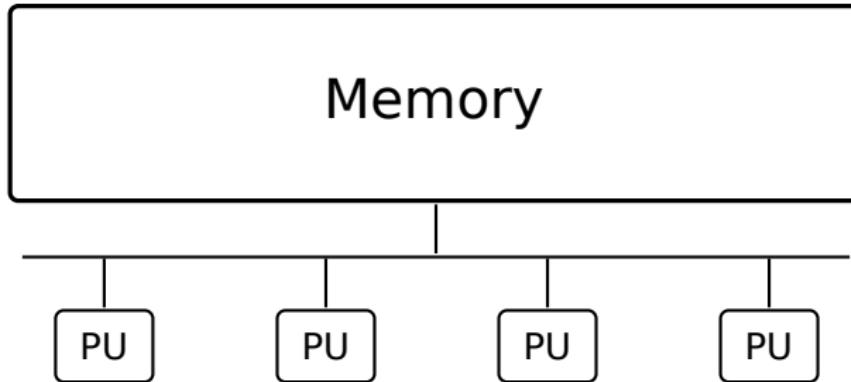
int main() {
    pthread_t tid;

    // A
    pthread_create(&tid, NULL, &mythread, NULL);

    // B
    pthread_join(tid, NULL);

    // D
}
```





- ▶ PThread
- ▶ OpenMP

OpenMP

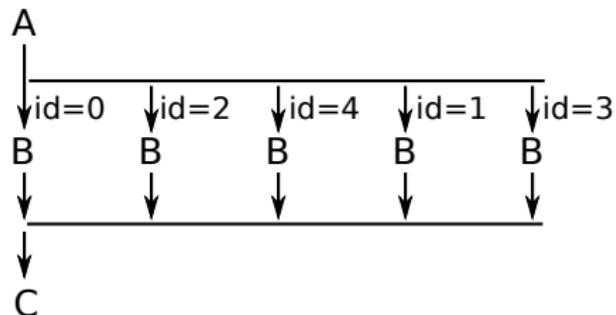
```
#include <omp.h>

int main() {
    // A

    #pragma omp parallel num_threads(5)
    {
        int id = omp_get_thread_num();
        printf("My id: %d\n", id); // B
    }

    // C
}
```

```
% gcc -fopenmp test.c
% ./a.out
My id: 1
My id: 3
My id: 0
My id: 2
My id: 4
```



OpenMP

```
#pragma omp parallel for num_threads(3)
for (int i=0; i<10; i++) {
    int id = omp_get_thread_num();
    printf("%d : %d\n", id, i);
}
```

```
% ./a.out
```

```
2: 7
```

```
2: 8
```

```
2: 9
```

```
1: 4
```

```
1: 5
```

```
1: 6
```

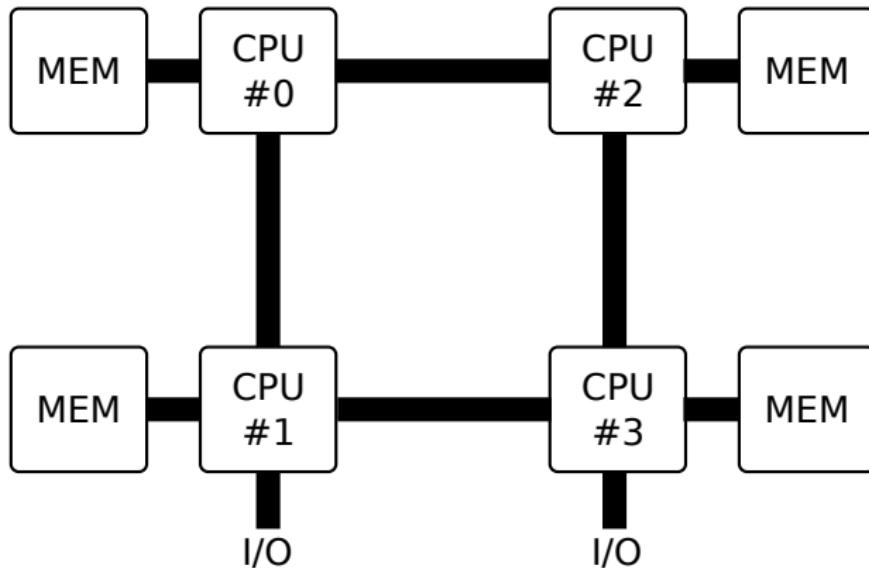
```
0: 0
```

```
0: 1
```

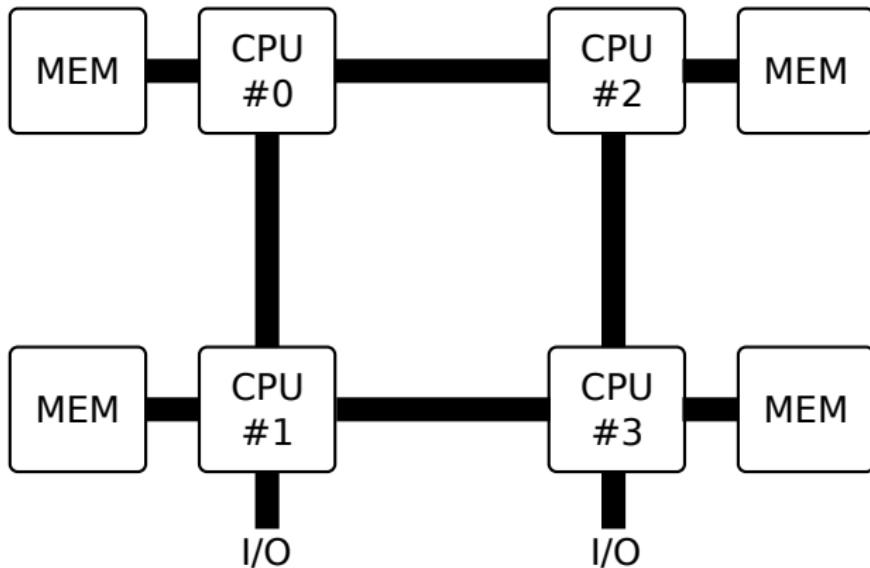
```
0: 2
```

```
0: 3
```

Multicore with NUMA

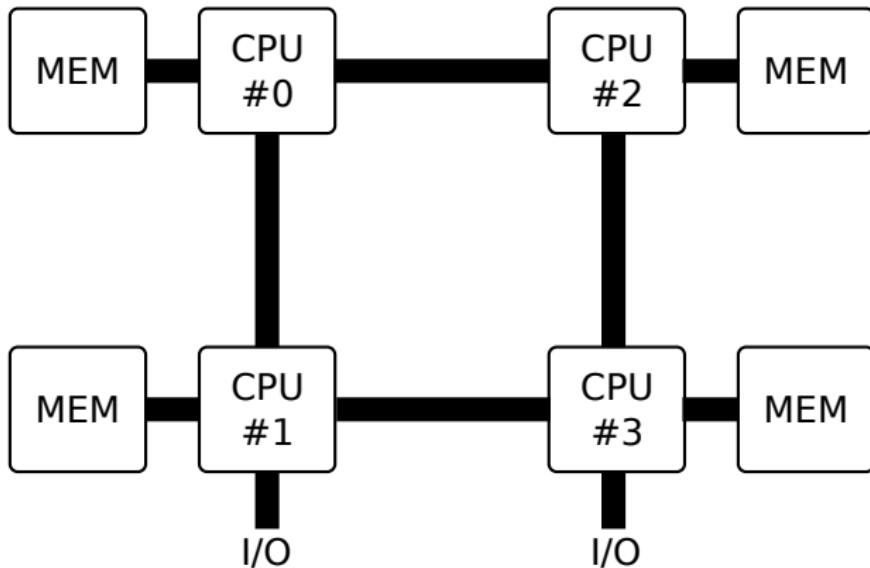


Multicore with NUMA



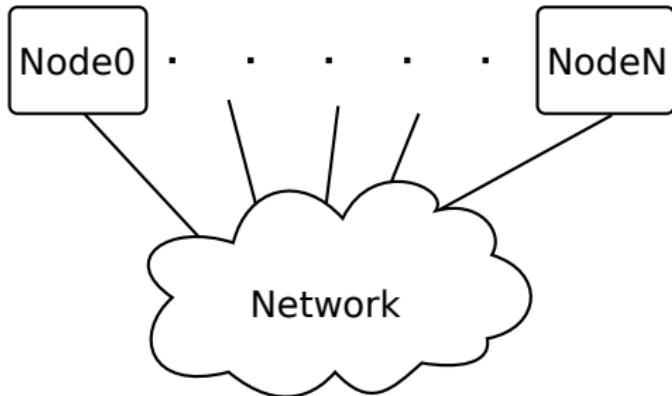
- ▶ Issues: thread scheduling (data affinity...), data migration...

Multicore with NUMA

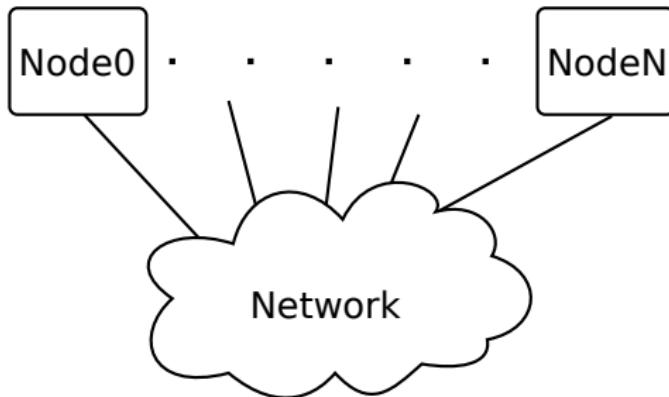


- ▶ Issues: thread scheduling (data affinity...), data migration...
- ▶ e.g. OpenMP: ForestGomp

Clusters



Clusters

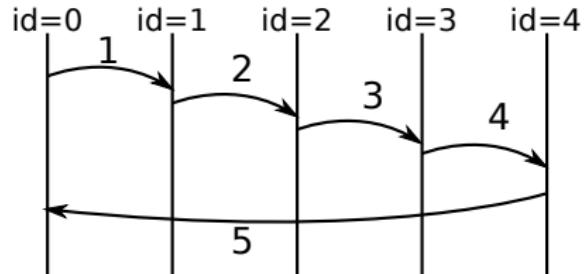


- ▶ Synchronous Message Passing: MPI

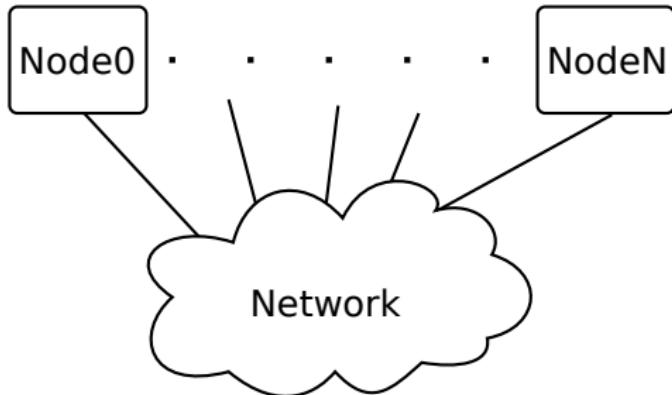
MPI

```
int cnt = 1;  
  
if (id == 0) {  
    MPI_Send(&cnt, 1, MPI_INT, id+1, 0, MPI_COMM_WORLD);  
    MPI_Recv(&cnt, 1, MPI_INT, N-1, 0, MPI_COMM_WORLD, NULL);  
    printf("Counter: %d\n", cnt);  
}  
else {  
    MPI_Recv(&cnt, 1, MPI_INT, id-1, 0, MPI_COMM_WORLD, NULL);  
    cnt += 1;  
    MPI_Send(&cnt, 1, MPI_INT, (id+1) % N, 0, MPI_COMM_WORLD);  
}
```

```
% mpicc test.c  
% mpirun -n 5 --hostfile hosts ./a.out  
Counter: 5
```



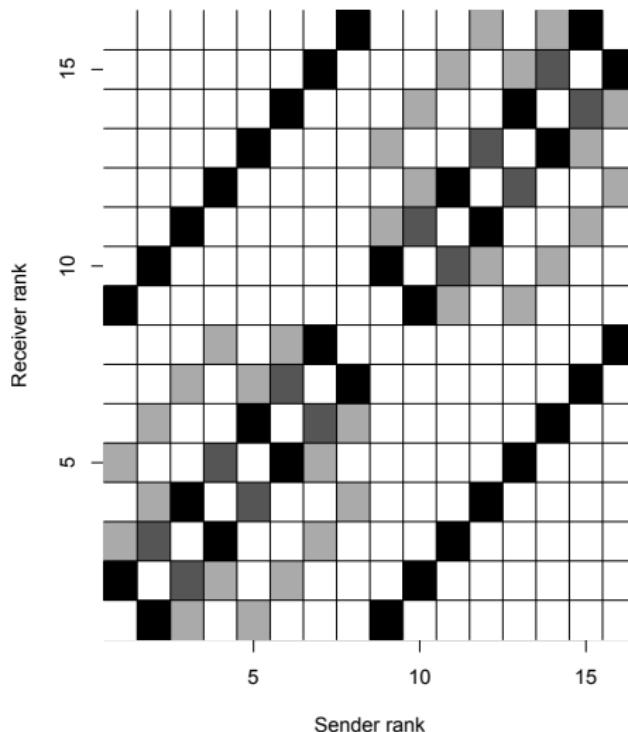
Clusters



- ▶ Synchronous Message Passing: MPI
 - ▶ Process placement issue
 - ▶ Depends on network topology
 - ▶ e.g. TreeMatch

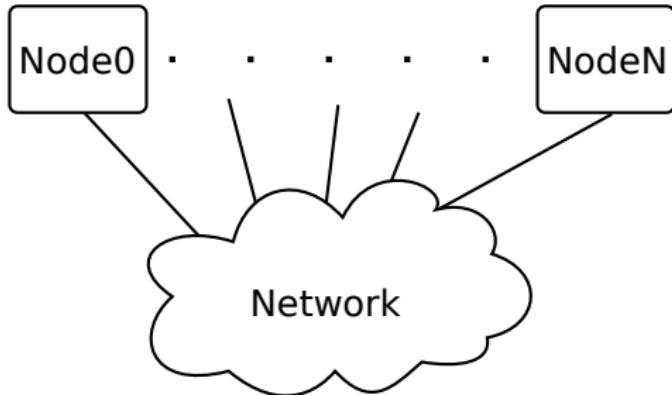
TreeMatch

ZEUSMP/2 communication pattern – Metric : msg



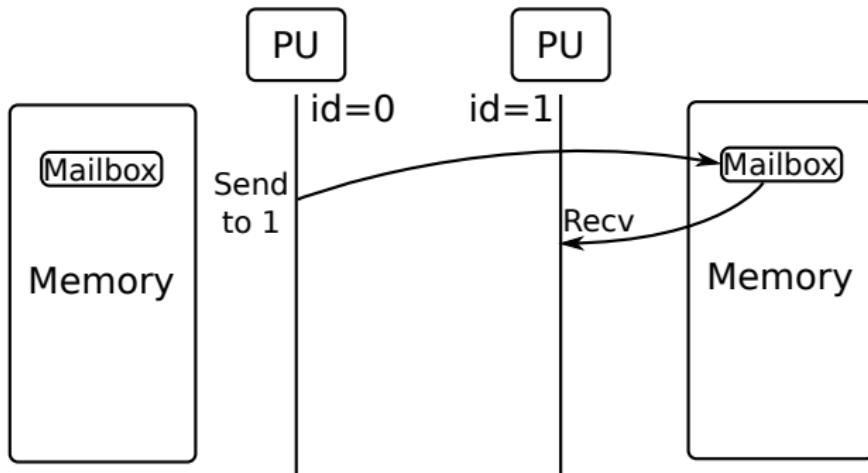
- ▶ Trace communications between processes
- ▶ Algorithm to map processes to network topology

Clusters



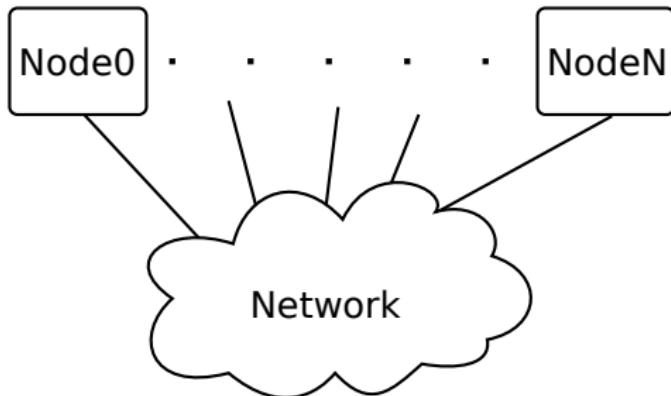
- ▶ Synchronous Message Passing: MPI
- ▶ Actor model / Active Objects: Charm++, Erlang...

Actor model



- ▶ Asynchronous message passing
- ▶ Dynamic actor creation

Clusters

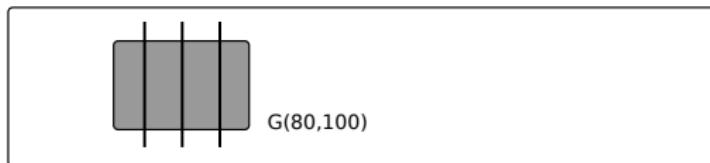


- ▶ Synchronous Message Passing: MPI
- ▶ Actor model / Active Objects: Charm++, Erlang...
- ▶ PGAS: HPF, XcalableMP, x10, ZPL/Chapel...
 - ▶ PGAS = Partitioned Global Address Space

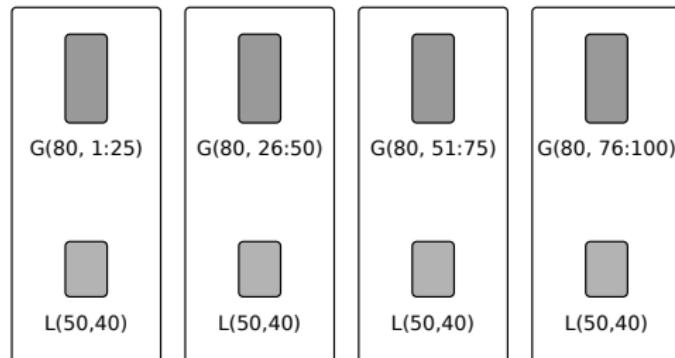
PGAS: XcalableMP

```
!$xmp nodes P(4)
!$xmp template T(100)
!$xmp distribute T(block) onto P
    real G(80, 100)           ! global variable
!$xmp align G(*, i) with T(i)
    real L(50, 40)            ! local variable (default)
```

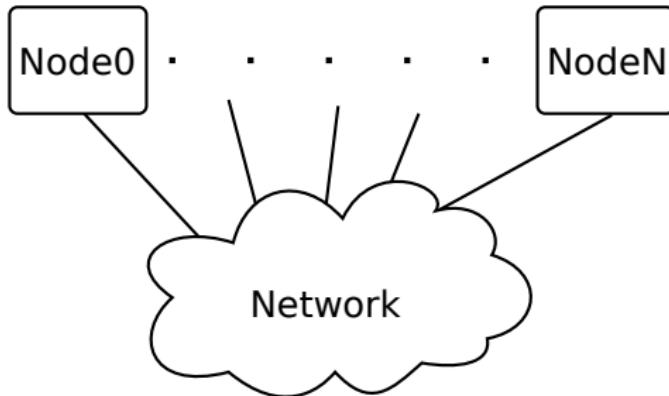
Global Address Space (virtual)



Local View (data allocation)

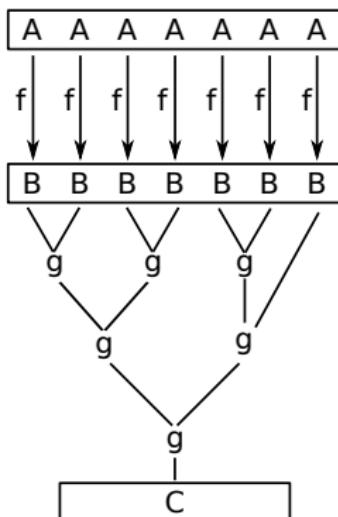


Clusters



- ▶ Synchronous Message Passing: MPI
- ▶ Actor model / Active Objects: Charm++, Erlang...
- ▶ PGAS: HPF, XcalableMP, x10, ZPL/Chapel...
- ▶ MapReduce: Google MapReduce, Hadoop...

MapReduce



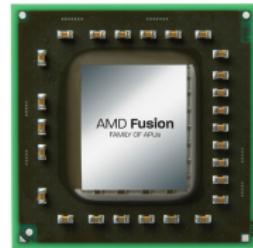
- ▶ reduce $g \circ map\ f$
 - ▶ Most operations on collections can be written this way
 - ▶ Easily parallelizable
- ▶ Users only have to provide f and g
- ▶ The runtime system provides
 - ▶ Task distribution
 - ▶ Communications
 - ▶ Fault tolerance
 - ▶ ...

Keywords: Bird-Meertens Formalism, Catamorphism

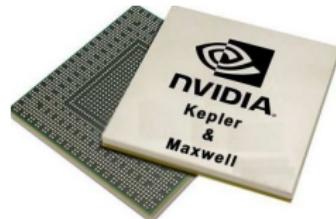
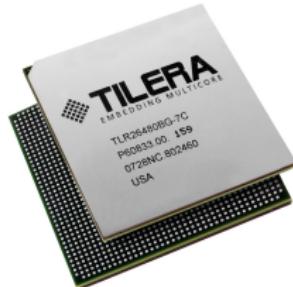
Summary

Architecture	Programming models
Multicore (NUMA)	OpenMP, PThread...
Cluster	MPI, PGAS, Actors, MapReduce...

- ▶ They all assume homogeneous architectures
 - ▶ Same (or similar) architecture duplicated



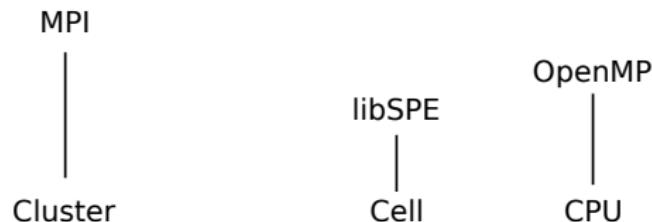
Accelerators & Many-Core



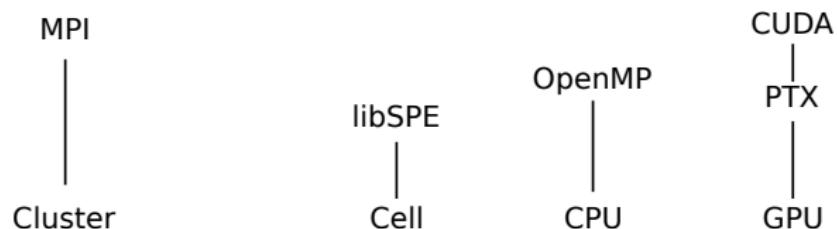
Frameworks . . .



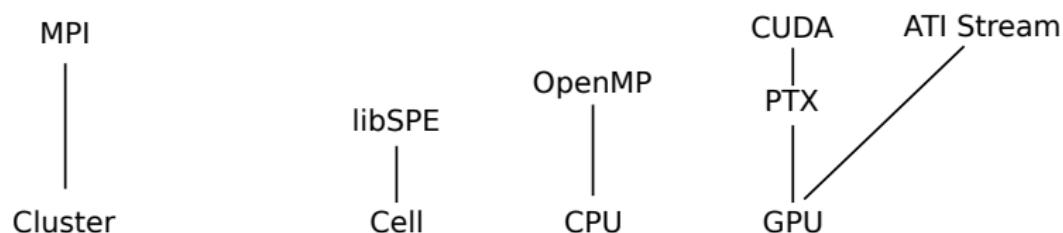
Frameworks...



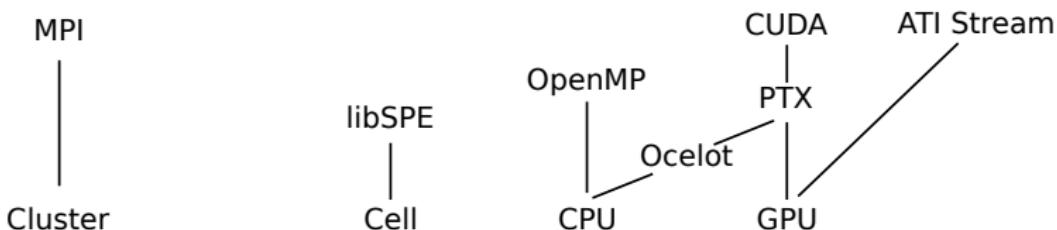
Frameworks...



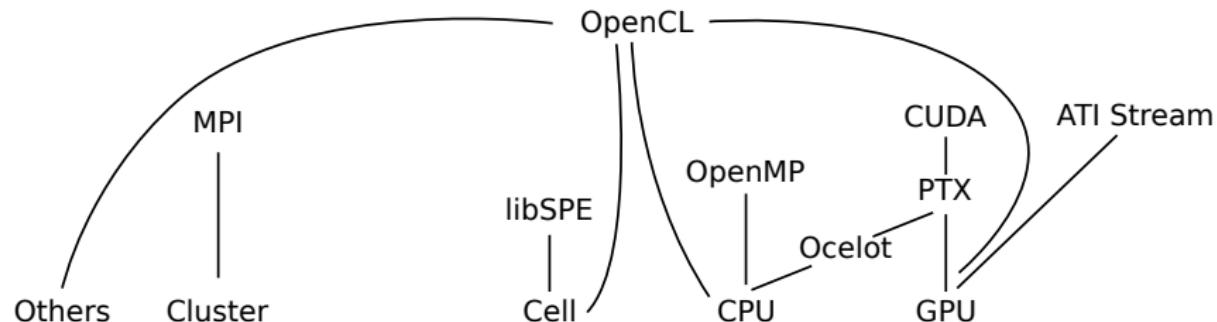
Frameworks...



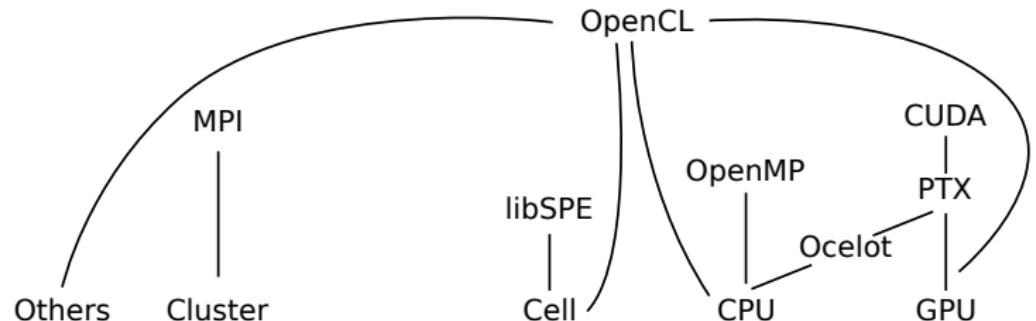
Frameworks...



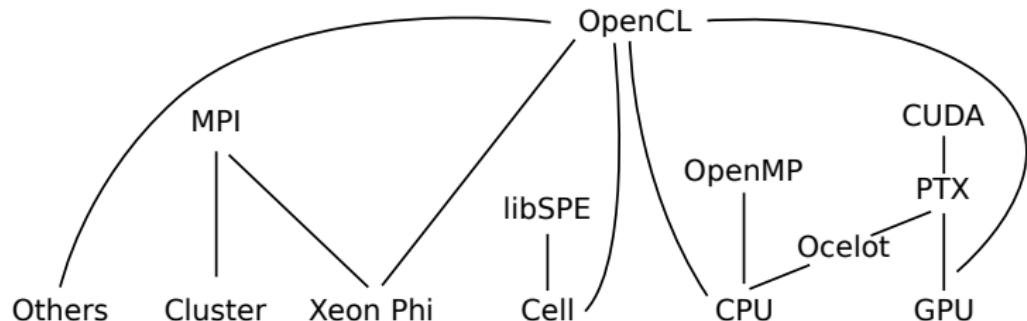
Frameworks...



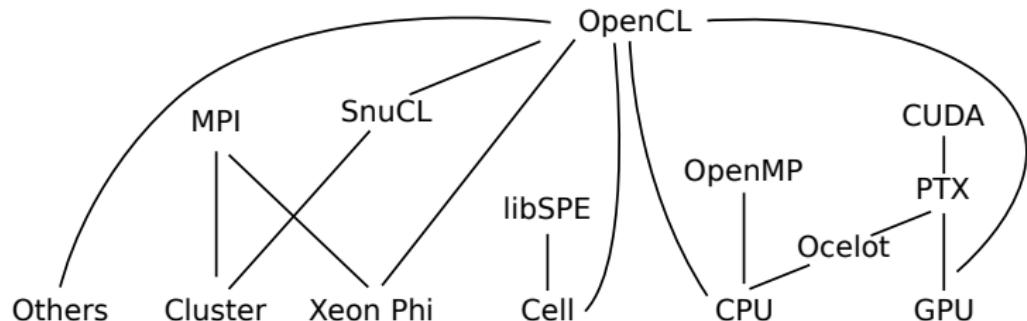
Frameworks...



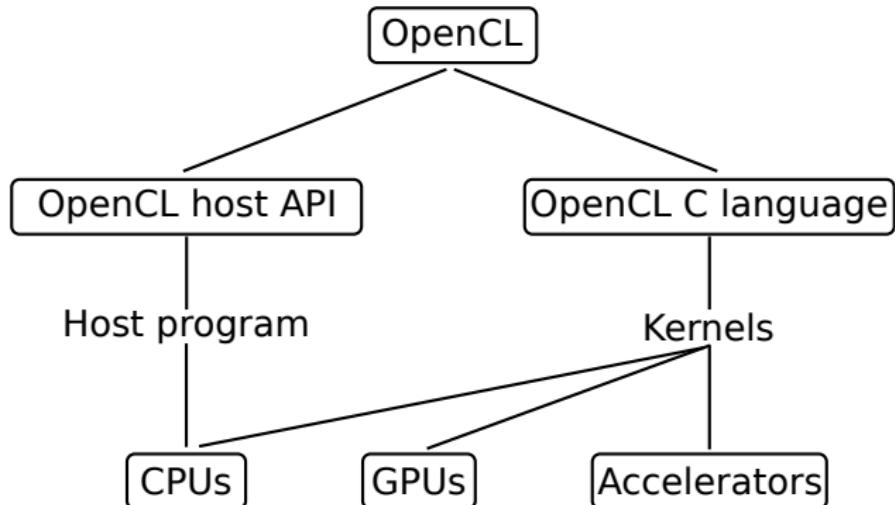
Frameworks...



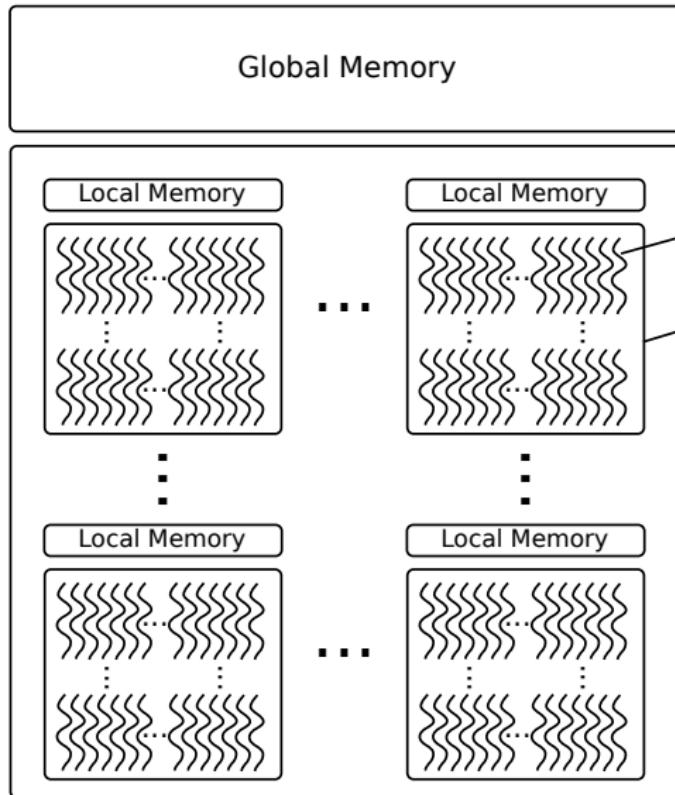
Frameworks...



OpenCL



OpenCL C Language



- ▶ SPMD
- ▶ `get_local_id(int dim)`
- ▶ `get_global_id(int dim)`

- ▶ `__local`, `__global`
- ▶ `work-group barrier`

OpenCL Host API

Device management

- ▶ Query available devices and their capabilities

Memory management

- ▶ Allocate/release buffers in device memory
- ▶ Transfer data between host memory and devices

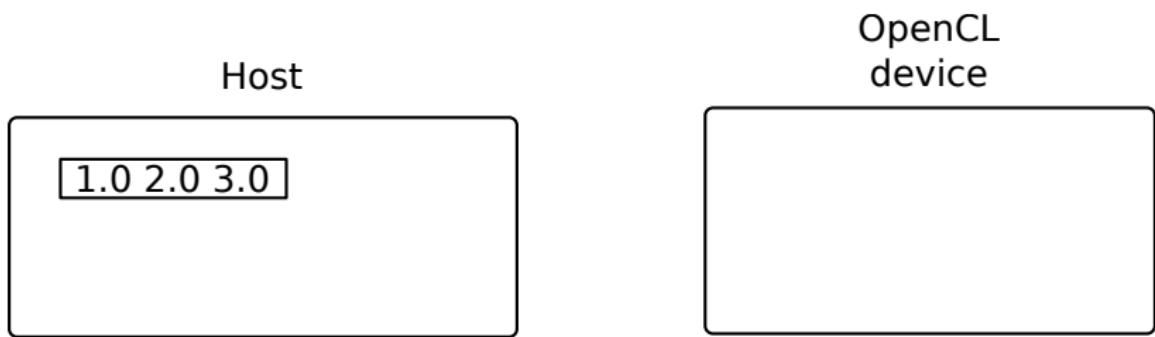
Kernel management

- ▶ Compilation
- ▶ Scheduling on devices

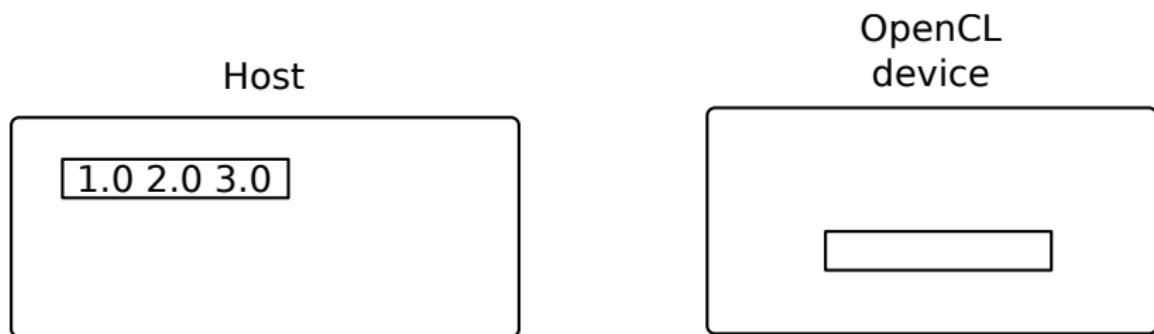
Synchronizations

- ▶ Command graph: events...
- ▶ With the host program: barriers...

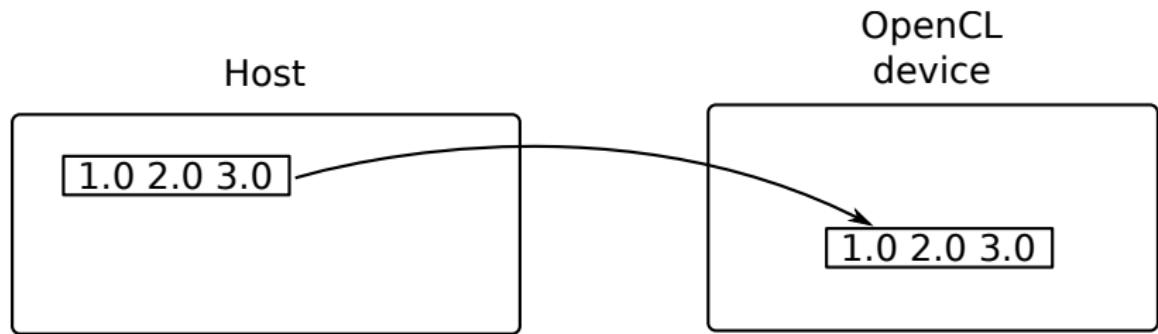
OpenCL example



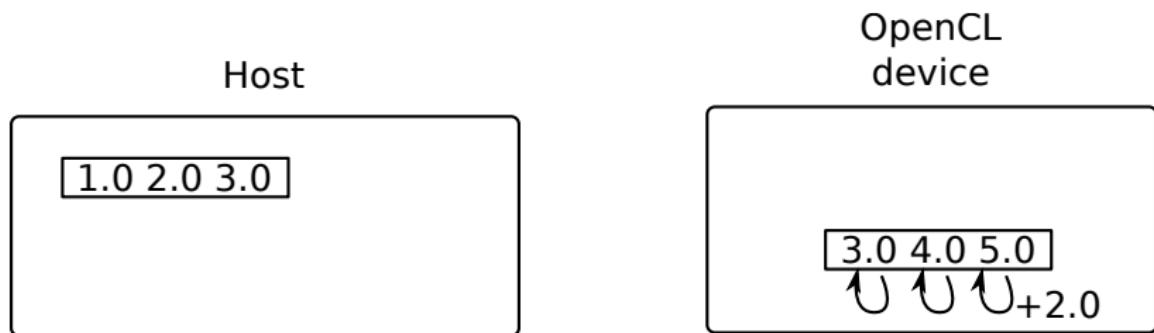
OpenCL example



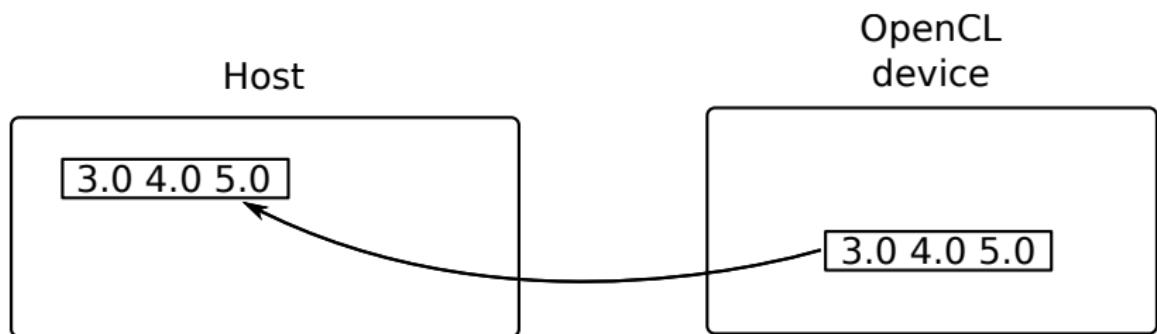
OpenCL example



OpenCL example



OpenCL example



OpenCL example: kernel

```
kernel void add2( __global float *out )
{
    size_t id = get_global_id(0);
    out[id] += 2.0;
}
```

OpenCL example: host API

```
float data[256] = { ... };

clGetPlatformIDs(max_pfs, pfs, NULL);
clGetDeviceIDs(pfs[0], CL_DEVICE_TYPE_ALL, 1, devs, NULL);
ctx = clCreateContext(NULL, 1, devs[0], NULL, NULL, NULL);
cq = clCreateCommandQueue(ctx, devs[0], 0, NULL);

buf = clCreateBuffer(ctx, CL_MEM_READ_WRITE, 1024, NULL, NULL);

clEnqueueWriteBuffer(cq, buf, CL_FALSE, 0, 1024, data, 0, NULL, NULL);

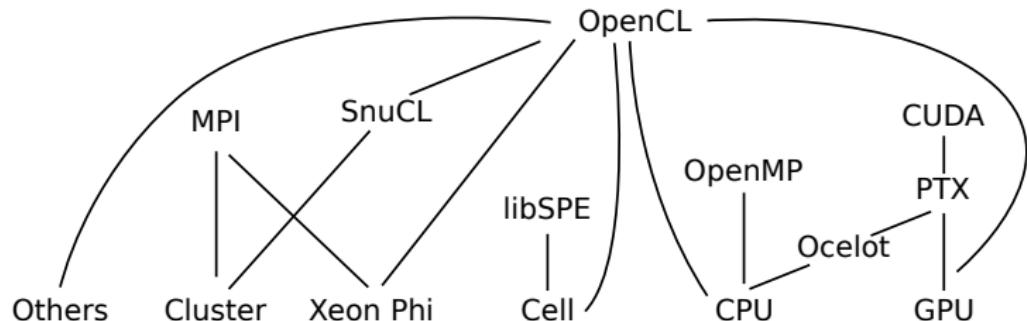
program = clCreateProgramWithSource(ctx, 1, &src, NULL, NULL);
clBuildProgram(program, 0, NULL, NULL, NULL);
ker = clCreateKernel(program, "add2", NULL);

clSetKernelArg(ker, 0, sizeof(cl_mem), &buf);
size_t group[3] = {32, 1, 1};
size_t grid[3] = {256, 1, 1};
clEnqueueNDRangeKernel(cq, ker, 3, NULL, grid, group, 0, NULL, NULL);

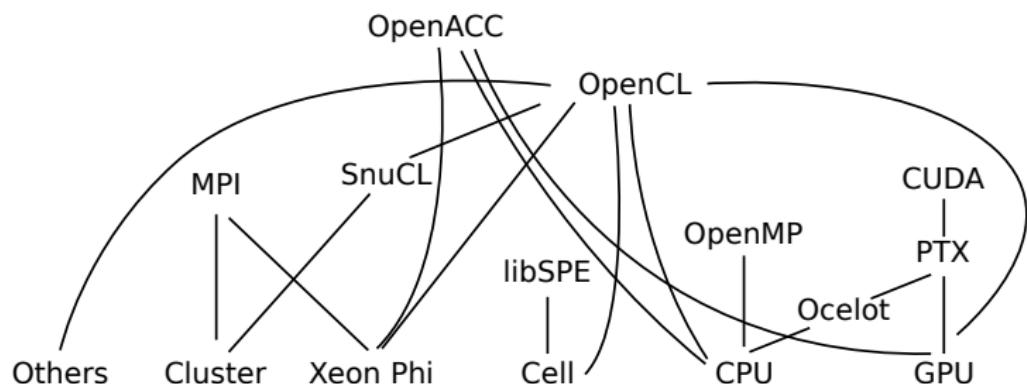
clEnqueueReadBuffer(cq, buf, CL_FALSE, 0, 1024, data, 0, NULL, NULL);

clFinish(cq);
```

Frameworks...



Frameworks...



OpenACC

Concepts

- ▶ Generate kernels from annotated C/Fortran code
- ▶ Simpler memory model (one accelerator)
- ▶ Ensure memory coherence with annotations

OpenACC example

```
float data[256] = { ... };

#pragma acc kernels copy(data)
for (int i=0; i<256; i++)
    data[i] += 2.0;
```

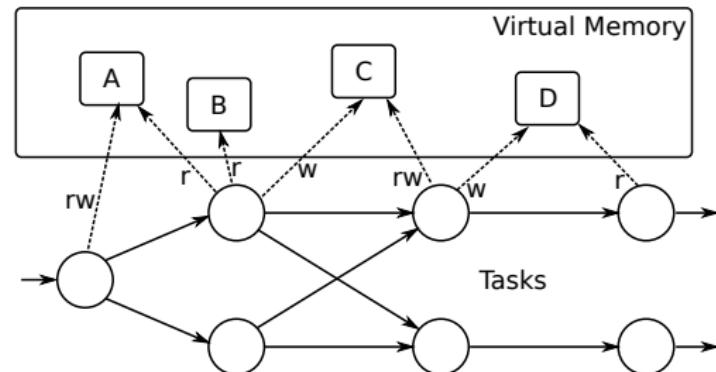
Limitations

Low-level approaches

- ▶ No support for multi-accelerator (OpenACC)
- ▶ Hard to manage several accelerators (OpenCL)

Runtime Systems

Runtime Systems



- ▶ Automatically schedule kernels on available accelerators
- ▶ Automatically manage distributed memory
- ▶ e.g. StarSS, StarPU...

Runtime System

Drivers (CPU, CUDA, OpenCL...)

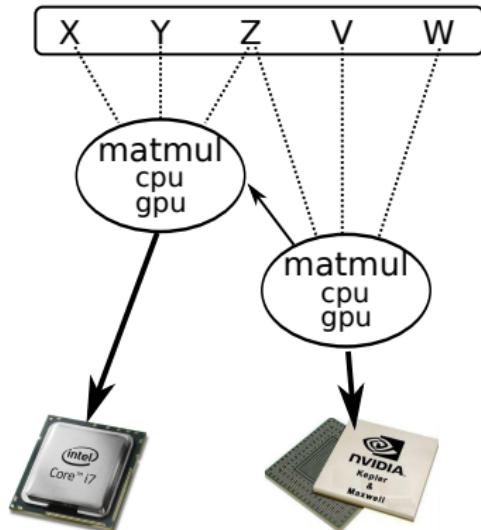
CPU

GPU

...

Xeon Phi

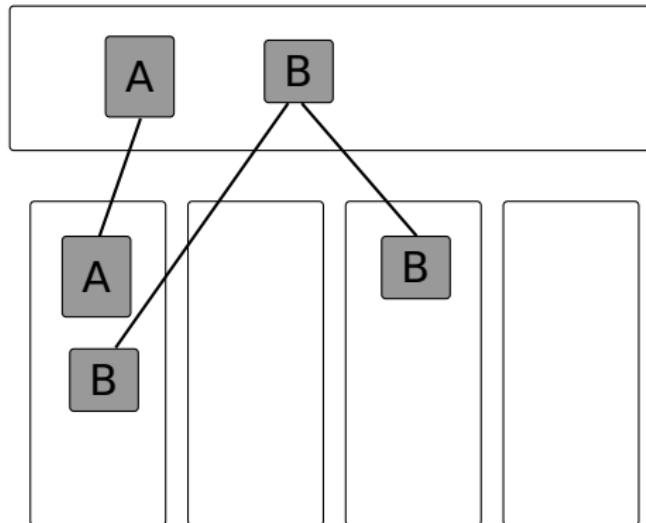
Runtime Systems



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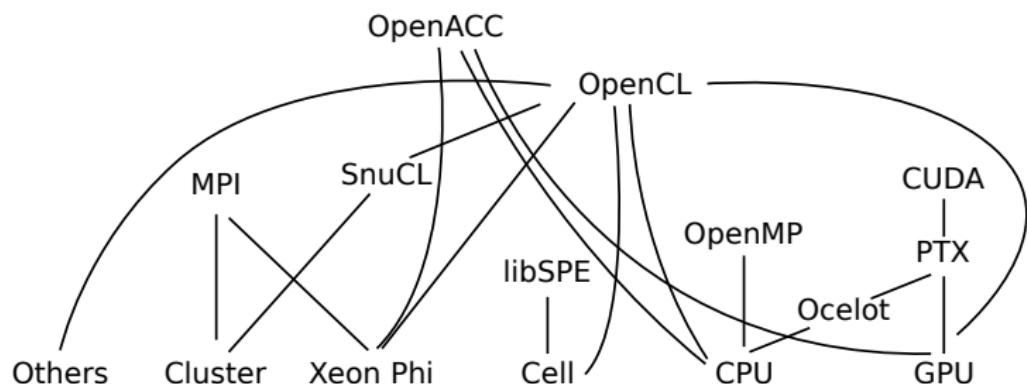
Runtime Systems

Global Space (virtual)

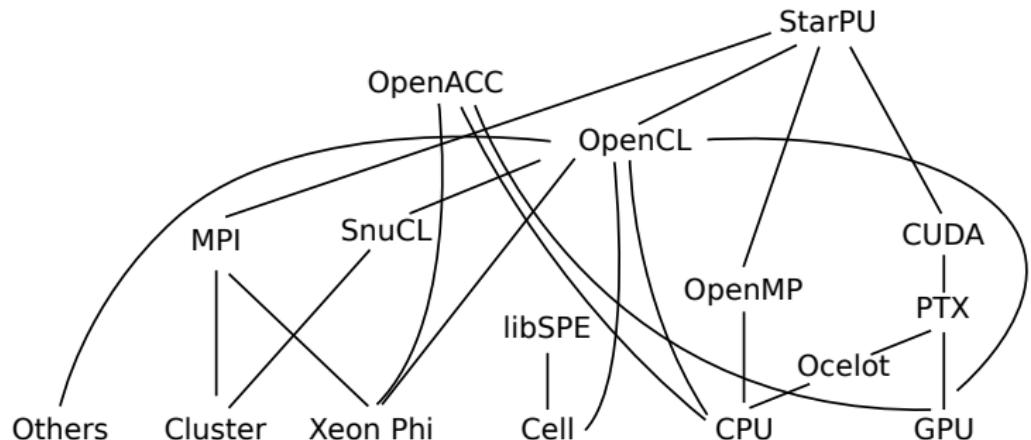


- ▶ Automatically schedule kernels on available accelerators
- ▶ Automatically manage distributed memory
- ▶ e.g. StarSS, StarPU...

Frameworks...



Frameworks...



StarPU example

```
void scal_cpu_func(void * buffers[], void *args);
void scal_cuda_func(void * buffers[], void *args);

starpu_codelet scal_cl = {
    .where = STARPU_CPU | STARPU_CUDA,
    .cpu_func = {scal_cpu_func, NULL},
    .cuda_func = {scal_cuda_func, NULL},
    .nbuffers = 1,
    .modes = {STARPU_RW}
}

float vector[N];

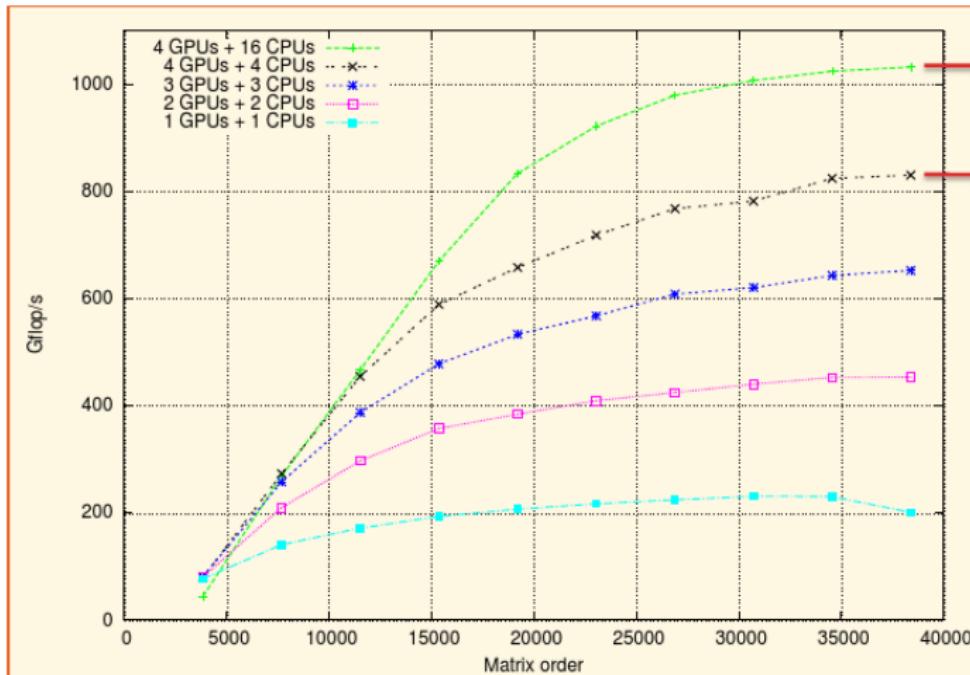
starpu_data_handle vector_handle;

starpu_vector_data_register(&vector_handle, 0,
                           vector, N, sizeof(float));

starpu_insert_task(&scal_cl, STARPU_RW, vector_handle, 0);

starpu_task_wait_for_all();
starpu_data_unregister(vector_handle);
```

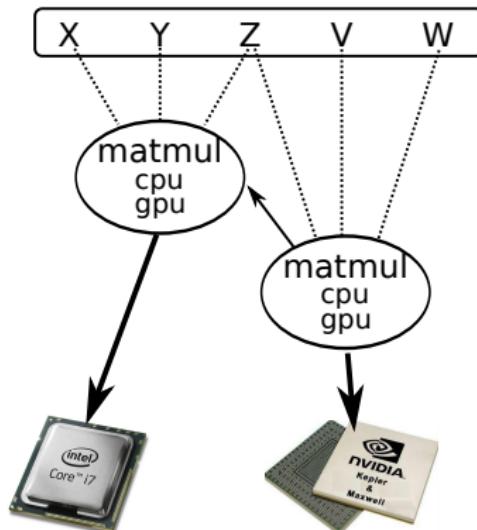
StarPU performance



UTK Magma (linear algebra) - QR factorization
16 CPUs (AMD) + 4 GPUs (C1060)

Limitations

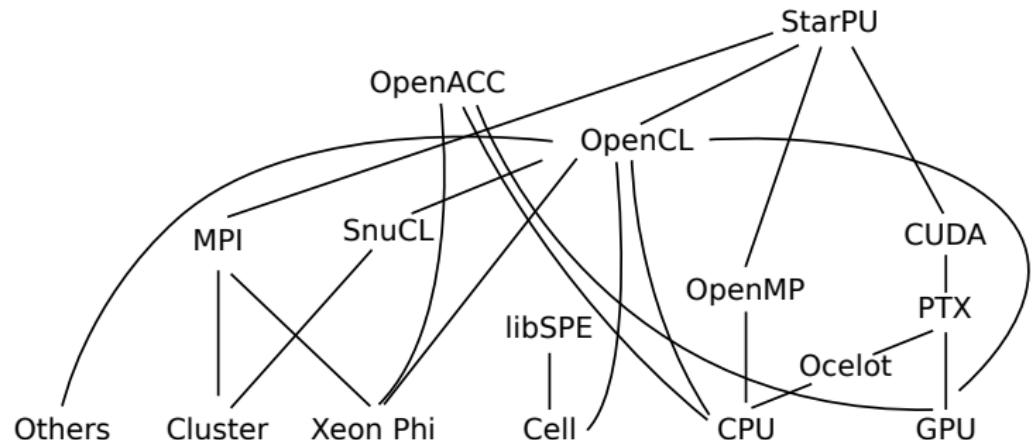
- ▶ Task/data granularity adaptation
- ▶ Task graph optimizations



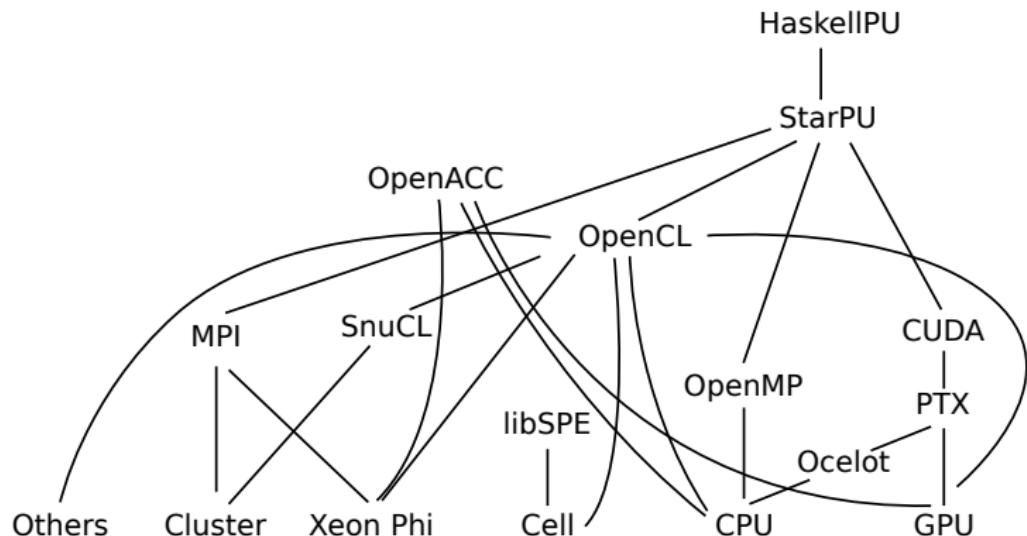
What I am working on

- ▶ Use a functional language to describe the task graph
- ▶ More amenable to optimizations
 - ▶ Granularity adaptation expected
- ▶ e.g. HaskellIPU, ViperVM

Frameworks...



Frameworks...

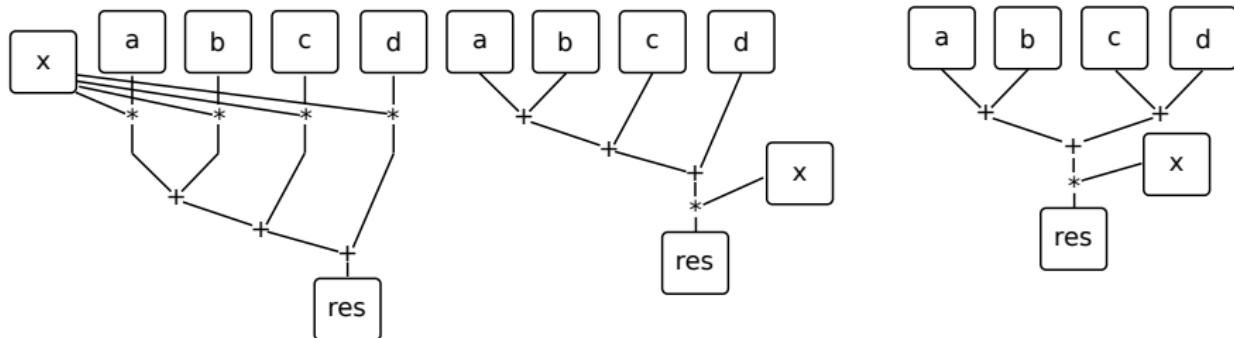


HaskellPU

```
let res = a*x + b*x + c*x + d*x  
computeSync res
```

```
let res = (a + b + c + d) * x  
computeSync res
```

```
let res = (reduce (+) [a,b,c,d]) * x  
computeSync res
```



Conclusion

- ▶ Things are evolving quickly!
 - ▶ Architectures
 - ▶ Programming models
- ▶ Runtime systems are now accepted in the HPC community
 - ▶ Ensure portability and performance
- ▶ New challenges
 - ▶ Fault tolerance
 - ▶ Power management
 - ▶ Collaboration with the programming language community (FP...)

Thank you!